DynamoDB Design Patterns

Intro to Dynamo

Dynamo is a managed, scalable, NoSQL key-value, wide-column database

managed: Dynamo exposes APIs and handles the rest

scalable: can handle up to 40k RPS; scales up and down with demand

NoSQL: not a relational database

key-value: think distributed hash table

wide-column: names / formats of item attributes varies from row-to-row

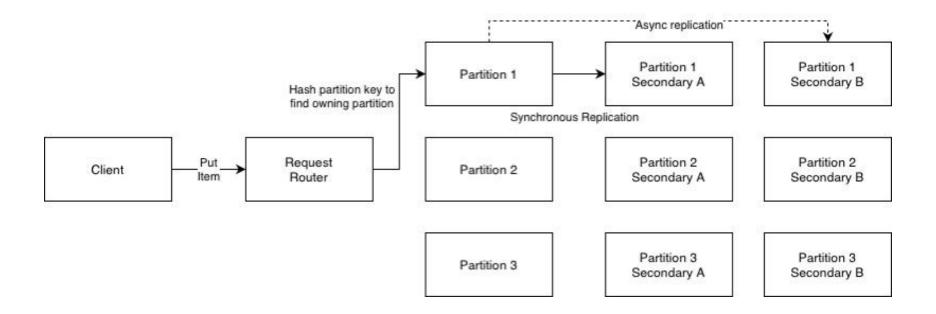
Why Dynamo?

- Highly available and scalable
- Strict query patterns guard against operations that won't scale
- Pay only for the capacity you need
- Tightly integrated with AWS (IAM, CloudWatch, CloudFormation)

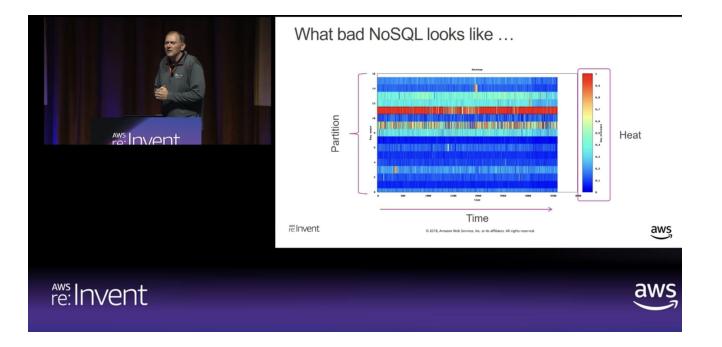
How It Works

- A Dynamo table is stored on a collection of nodes (partitions)
- Each table has a **partition key** and optional **sort key**, which uniquely identify all items in table
- **Partition key** decides which node owns the record
- **Sort key** determines how records are organized within a node
- Each item is collection of attributes, which can be scalar (string, number, binary, boolean) or complex types (list, map, set)

How It Works



Choosing Partition and Sort Keys



Choosing Partition and Sort Keys

- Key selection drives access patterns
- Partition key must be **high cardinality** to avoid hot partitions
 - Good: GUID, CustomerId
 - Bad: Status, Boolean
- Sort key can be used for ordering and modeling 1:n and n:n relationships
 - Relationships modeled with composite keys
 - Order can be maintained with timestamps (updated_at) or sortable GUIDs, e.g. KSUID

Indexing

- Global secondary indexes (GSI) allow you to define a new partition key and sort key on the table
 - Enables new "views" on a table
 - RCU/WCU must be at least equal to table, or throttling may happen
 - GSI are only eventually consistent
- Local secondary indexes (LSI) allow you to define a new sort key on existing table partition key
 - Reorganizes data in a single partition. Imposes a 10 GB limit per hash key
 - Strongly consistent, as opposed to GSI

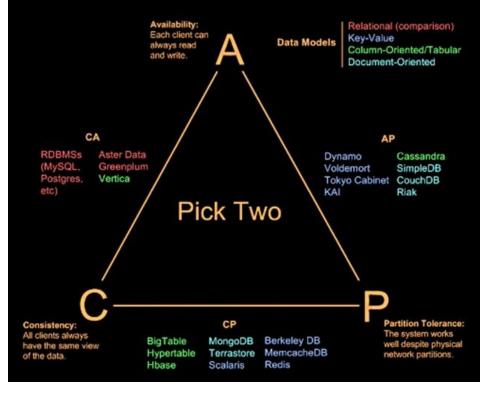
Consistency

- Writes always go to the owning node and are consistent
- Eventually consistent reads go to any partition
- Strongly consistent reads go to owning partition and cost 2x
- Prefer eventual consistency when possible
- **Only** table partition key / sort key and local secondary indexes can provide full read consistency
- TransactionWriteItems / TransactionGetItems can provide ACID compliance

Dynamo & CAP Theorem

- All distributed data stores can only provide 2 of 3
 - Consistency
 - Availability
 - Partition Tolerance
- Dynamo by default provides availability and partition tolerance
- Using strong consistency/transactions trades availability for consistency

Visual Guide to NoSQL Systems



Streams

- Provides a time-ordered sequence of item-level changes on a table
- Backed by Kinesis stream; out of the path of table requests
- Great for extending a Dynamo table with reactive functionality

Note on Capacity

- Billed by read capacity units (RCU) and write capacity units (WCU)
- Either on demand or provisioned modes
- Tip: Until your Dynamo workload is known, use on demand
- Cost of over-provisioning will likely exceed on demand costs

Limits

- Up to 40,000 RCU/WCU per table
- Max partition key size 2048 bytes
- Max sort key size is 1024 bytes
- Max 400KB item size
- Only 1MB of table data scanned per query before filters applied
- Single partition can only have 3000 RCU / 1000 WCU
 - \circ In other words, a key item cannot be written to > 1000 WCU
- Local secondary index can only contain 10GB of data per partition key
- 20 GSIs per table
- 5 LSIs per table

Single Table Design

- Different entities can & should live in the same Dynamo table
 - 1 table per entity (e.g. Users table, Roles table, Org table) is often an antipattern
 - Attempting to join across Dynamo tables can kill performance
- Solution: "pre-join" records of different types into a single table
- Partition keys allow us to specify "item collections"
- Sort keys allow us to define relationships between items
- GSIs / LSIs give us additional querying flexibility

Dynamo Patterns

Single Table Design Example

- Usecase: modeling simple e-commerce site
- Entities: User, Order, Product, Inventory

Primary Key					lhutaa					
PK	SK		Attributes							
	#PROFILE#alexdebrie	Username	FullName	Email	CreatedAt	Addresses				
		alexdebrie	Alex DeBrie	alexdebrie1@gmail.com	03/23/2018	{"Home":{"StreetAddress":"1111 1st St","State":"Nebr				
	ORDER#5e7272b7	Username	Orderld	Status	CreatedAt	Address				
LICE D#alavdabria		alexdebrie	5e7272b7	PLACED	04/21/2019	{"StreetAddress":"1111 1st St","State":"Nebraska","Co				
USER#alexdebrie	ORDER#42ef295e	Username	Orderld	Status	CreatedAt	Address				
		alexdebrie	42ef295e	PLACED	04/25/2019	{"StreetAddress":"1111 1st St","State":"Nebraska","Co				
	ORDER#2e7abecc	Username	Orderld	Status	CreatedAt	Address				
		alexdebrie	2e7abecc	SHIPPED	12/25/2018	{"StreetAddress":"1111 1st St","State":"Nebraska","Co				
	#PROFILE#nedstark	Username	FullName	Email	CreatedAt	Addresses				
		nedstark	Eddard Stark	lord@winterfell.com	02/27/2016	{"Home":{"StreetAddress":"1234 2nd Ave","City":"Wir				
	ORDER#2eae1dee	Username	Orderld	Status	CreatedAt	Address				
USER#nedstark		nedstark	2eae1dee	SHIPPED	01/15/2019	{"StreetAddress":"Suite 200, Red Keep","City":"King's L				
		Username	Orderld	Status	CreatedAt	Address				
	ORDER#f4f80a91	nedstark	f4f80a91	PLACED	05/12/2019	{"StreetAddress":"Suite 200, Red Keep","City":"King's L				

Sca	an [Table]	AmazonExample: PK, SK		V	~
	 Add filte 				
	Start sear	ch			
	РК	SK -	Address ()	Description	Name
	USER#will	USER#will	{"office": "DCA15", "address": "17		
	USER#will	ORDER#1			
	USER#will	ORDER#2			
	PRODUCT#1	INVENTORY#DCA			
	PRODUCT#1	INVENTORY#SEA			
	PRODUCT#1	PRODUCT#1		Office Chair	Chair
	USER#jdoe	ORDER#1			
	USER#jdoe	USER#jdoe			

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Query: [Table	e] AmazonExa	mple: PK, SK 🖍	•					Viewing	g 1 to 1 item
Query [Table] AmazonExample: PK, SK									
Partition key	PK		String				USER#will		
Sort key	SK		String		=		USER#will		
	Add filter								-
Sort	Ascending	g 🔵 Descending							
Attributes	💿 All 🔵 Pro	ojected							
	Start search	1							-
РК	Ŧ	SK	-	Addres	SS				-
USER#	will	USER#will		{"office	e": "DCA15",	addres"	ss": "1775 Belle St", "ci	ty": "Arlington", "s	state":

Query: [Table] AmazonExample: PK, SK



Query V [Table] AmazonExample: PK, SK										
Partition key	PK	Strin	g			USER#will				
Sort key	SK	Strin	g	=		Enter value				
	Add filter									
Sort	Sort 💿 Ascending 🔘 Descending									
Attributes	Attributes All Projected									
Start search										
РК	*	SK -	Addres	s ()		Items				
USER#	ŧwill	USER#will	{"office"	: "DCA15", "ad.						
USER#	ŧwill	ORDER#1				[{ "M" : { "Id" : { "S" : "PRODUCT	「#1" }, "Name" : {			
USER#	ŧwill	ORDER#2				[{ "M" : { "Id" : { "S" : "PRODUCT	「#1" }, "Name" : {			

Query: [Table] AmazonExample: PK, SK



Query V [Table] AmazonExample: PK, SK										
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	Add filter									
Sort	Sort Ascending Descending 									
Attributes	Attributes 💿 All 🔵 Projected									
	Start search									
РК	Ŧ	SK	-	Description	Ŧ	Name	Ŧ	Quantity -		
PRODU	JCT#1	INVENTORY#DC	A					10		
PRODU	JCT#1	INVENTORY#SE/	Ą					2		
PRODU	JCT#1	PRODUCT#1		Office Chair		Chair				

Single Table Design Tips

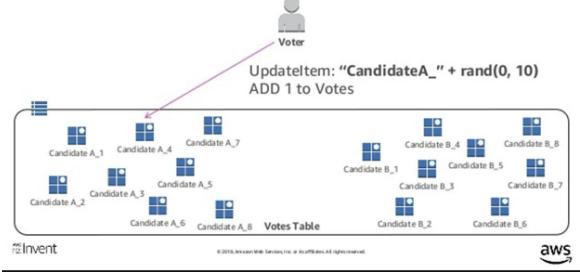
- Define access patterns up-front
- Don't lock partition key and sort key into one usecase
- Attempt to de-normalize data where possible to reduce number of queries
- Leverage sortable IDs to maintain order
- Try to keep partition and sort key identifiers short to prevent hitting size limits

Uniqueness Constraints on Multiple Attributes

- Uniqueness constraints can be added with new entity types
- Example: enforcing uniqueness on both email and user ID attribute
- Create write transaction
 - Item 1 {PK: "USER#3921", SK: "USER#3921"}
 - Item 2 {PK: "EMAIL#will@gmail.com", SK: "EMAIL#will@gmail.com"}
 - Conditional check on neither existing
- If either exists, the transaction fails

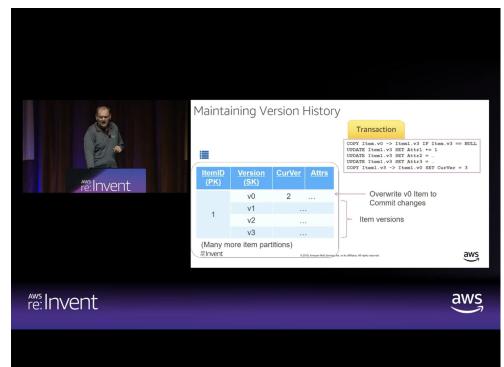
Avoiding Hot Keys

- Example: writing many votes to a candidate record
- To avoid high WCU to one item, shard the item among many records, and compute an aggregate count via an async process

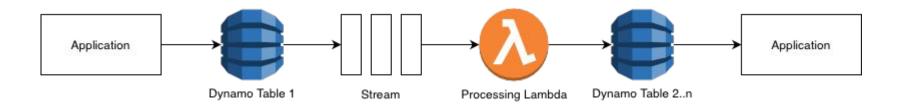


Maintaining Item History

- Using version tag as sort key allows maintaining of write history
- Use the following write pattern:
 - Fetch item ID = X, version = v0
 - In write transaction,
 - Set previous v0 as new vX item
 - Update attributes on v0
- v0 always contains latest record

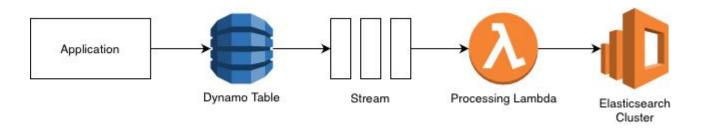


Aggregation with Streams

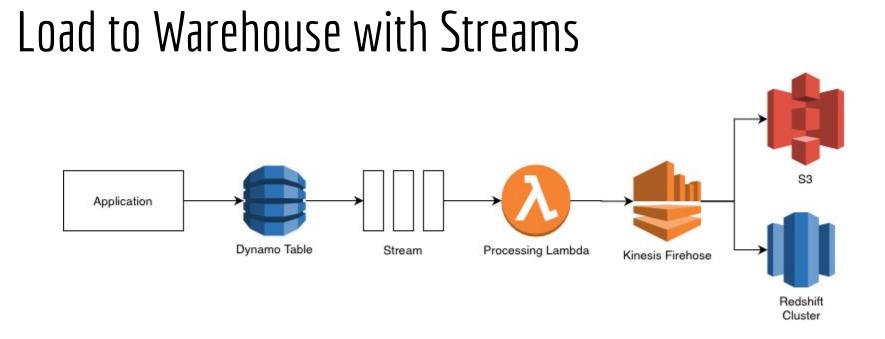


- Read updates from stream, and push metadata / aggregations back to Dynamo
- Example: on new Order item, update User record openOrders += 1

Full Text Search with Streams

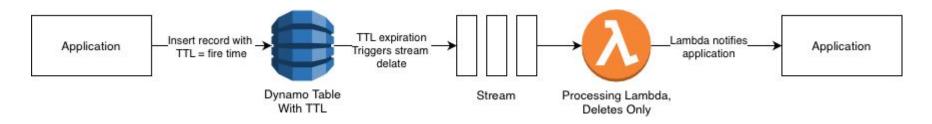


- Transform and push items into ElasticSearch to enable full-text search
- Maintains Dynamo as source of truth, but enables more powerful querying options



• Push items into a data warehouse (e.g. S3, Redshift) to enable flexible BI querying

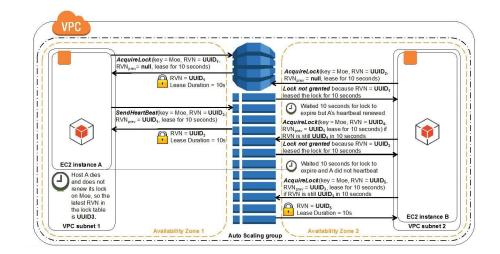
Scheduling with Item TTL



- Documents in Dynamo can define a time-to-live attribute. This is helpful for caching, leases, and scheduling
- Can schedule events by pushing an item with a TTL at desired fire date
- Listen for delete events, and notify application
- TTL is implemented as a background scanning process on each partition. Depending on table usage, delete could take as long as 48h to process. Typically is much shorter

Distributed Locking

- Quickly build distributed lock with open-source Amazon DynamoDB Lock Client
- Use cases: making sure two workers don't operate on same entity, leader election
- Only requires a Dynamo DB table with partition key "key"
- Supports heartbeats, lease duration, blocking / non-blocking lock acquisitions



Optimistic Locking

- Java library has @DynamoDBVersionAttribute annotation, which sets up optimistic item locking
- Each item put is given a conditional check, where current version = expected version
- Write will reject if versions do not match
- Make sure to handle ConditionalCheckFailed runtime exceptions

Serverless Dynamo Frontends

- AWS API Gateway can add REST endpoints on top of Dynamo tables using service proxies
 - No-code solution
 - APIs can be versioned
 - SIGv4 Authentication
 - E.g. map route GET /companies/Amazon/employees/1 to query PK: Amazon, SK: 1; transform and return result
- AWS AppSync can add GraphQL operations on top of Dynamo tables

From Millisecond to Microsecond

- Dynamo Accelerator (DAX) is fully managed caching solution which brings latencies down to microseconds
 - In-memory cache of items and queries
 - Only supports eventually consistent reads
- Global tables can replicate Dynamo tables cross-region
 - Multi-master replicas
 - Writes propagated cross-region within a second
 - Last-writer-wins for cross-region write conflicts

Handling Migrations

- If possible, handle new attribute defaults in business logic
- For small backfills / migrations, scripts are preferable
 - Parallel scans possible with TotalSegments / SegmentNumber arguments
- For large backfills / migrations, use EMR
 - AWS EMR has built in Dynamo adapters
 - Load a Dynamo table into Hive, make change, then load back to Dynamo
 - Make sure to set a % of capacity to use during job

Resources

- ReInvent 2018 DAT401 Advanced Design Patterns for DynamoDB https://youtu.be/HaEPXoXVf2k
- The DynamoDB Book by Alex Debrie https://www.dynamodbbook.com
- Advanced Design Patterns for Amazon DynamoDB by National Australia Bank https://link.medium.com/ypcCdKt6Kbb
- Dynamo Docs https://docs.aws.amazon.com/dynamodb/index.html